

Introduction to Computer Science: Programming Methodology

Lecture 10 Linked List

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Why we need another list data type

 Python's list class is highly optimized, and often a great choice for storage

 However, many programming languages do not support this kind of optimized list data type

More about linked lists



Common Questions

When are linked lists preferred over lists?

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- I'm learning data structures and algorithms using Python. I've learnt
 that the advantage of linked list is that it does not have a maximum number of nodes, unlike arrays in other languages.
 - Since Python automatically resizes our lists for us, the advantage has been abstracted away for us.
 - Therefore, I've always thought that the only advantage linked lists have was that adding a node at the front or back of the linked list was O(1), whereas adding an element to a list could end up being O(n) from Python having to resize the array and copving every element over.

Why do we need to build linked lists in Python?



Common Questions

- Linked list is a **classic** data structure as a good example.
- Rarely used in **high** level programming languages like Python. As

a practitioner, you may never use it in Python.

• We use Python to demonstrate the **concept**, so you know what it

is when using non-Python languages.

List in Python is a referential structure

List in Python is a referential structure





List in Python is a referential structure



Compact array

- A collection of numbers are usually stored as a compact array in languages such as C/C++ and Java
- A compact array is storing the bits that represent the primary data (not reference)
- The overall memory usage will be much lower for a compact structure because there is no overhead devoted to the explicit storage of the sequence of memory references (in addition to the primary data)

Linked List

- A singly linked list, in its simplest form, is a collection of nodes that collectively form a linear sequence
- Each node stores a reference to an object that is an element of the sequence, as well as a reference to the next node of the list



Linked List

- The first and last nodes of a linked list are known as the head and tail of the list, respectively
- By starting at the head, and moving from one node to another by following each node's next reference, we can reach the tail of the list
- We can identify the tail as the node having None as its next reference. This process is commonly known as traversing the linked list.
- Because the next reference of a node can be viewed as a link or pointer to another node, the process of traversing a list is also known as link hopping or pointer hopping





class Node:

def __init__(self,element,pointer=None):
 self.element = element
 self.pointer = pointer

head = Node(0)
node1 = Node(1)
node2 = Node(2)
tail = Node(3)

head.pointer = node1 node1.pointer = node2 node2.pointer = tail

p = head while(p!=None): print(p.element) p = p.pointer

Inserting an Element at the Head of a Singly Linked List







Pseudo code for inserting a node at the head

Algorithm add_first(L,e):

Inserting an Element at the Tail of a Singly Linked List



Pseudo code for inserting at the tail

```
Algorithm add_last(L,e):
```

Linked list :: insert

Add a new element to the list



Removing an Element from the head of a Singly Linked List



Pseudo code for removing a node from the head

Algorithm remove_first(L):

- if L.head is None then
 - Indicate an error: the list is empty.
 - L.head = L.head.next
 - L.size = L.size 1

{make head point to next node (or None)} {decrement the node count}

Linked list :: delete

• Delete a node from the list



Practice: Implement stack with a singly linked list

```
class Node:
    def __init__(self, element, pointer):
        self.element = element
        self.pointer = pointer
```

class LinkedStack:

```
def __init__(self):
    self.head = None
    self.size = 0
def __len__(self):
    return self.size
def is_empty(self):
    return self.size == 0
def push(self, e):
    self.head = Node(e, self.head)
    self.size += 1
```

```
def top(self):
    if self.is_empty():
        print('Stack is empty.')
    else:
        return self. head. element
def pop(self):
    if self.is_empty():
        print('Stack is empty.')
    else:
        answer = self. head. element
        self.head = self.head.pointer
        self. size -=1
        return answer
```

Practice: Implement queue with a singly linked list

class LinkedQueue:

```
def __init__(self):
    self. head = None
    self. tail = None
    self. size = 0
def __len__(self):
    return self. size
def is_empty(self):
    return self. size == 0
def first(self):
    if self.is_empty():
        print('Queue is empty.')
    else:
        return self, head, element
```

```
def dequeue(self):
    if self.is_empty():
        print('Queue is empty.')
    else:
        answer = self. head. element
        self. head = self. head. pointer
        self.size -= 1
        if self.is empty():
            self.tail = None
        return answer
def enqueue(self, e):
    newest = Node(e, None)
    if self.is empty():
        self.head = newest
    else:
        self.tail.pointer = newest
    self.tail = newest
    self. size += 1
```

Practice: how to find the middle node (Optional)



Practice: how to find the middle node (Optional)

```
def findMiddleByLength(head):
    # First loop to calculate the length of the linked list
    length = 0
    current = head
    while current:
        length += 1
        current = current.next
```

```
# Calculate the middle index for the first middle node
middle_index = (length - 1) // 2
```

```
# Second loop to traverse to the middle node
current = head
index = 0
while index < middle_index:
    current = current.next
    index += 1</pre>
```

```
print("The middle node's data is:", current.data)
return current
```

Practice: how to find the middle node (Optional)

```
# Define the Node class
class Node:
    def __init__(self, data):
        self.data = data # Store data
        self.next = None # Pointer to the next node
# Function to find the middle node of the linked list
def findMiddle(head):
    slow = head # Slow pointer moves one step at a time
    fast = head # Fast pointer moves two steps at a time
    while fast and fast.next:
        slow = slow.next # Move slow pointer by one
        fast = fast.next.next # Move fast pointer by two
    print("The middle node's data is:", slow.data)
    return slow
```

How to reverse a linked list (Optional)



New head

How to reverse a linked list (Optional)

```
class Node:
   def __init__(self, data):
       self.data = data # Store data
       self_next = None # Pointer to the next node
# Function to reverse the linked list
def reverseLinkedList(head):
   prev = None
   current = head
   while current:
       next_node = current.next  # Save the next node
                           # Reverse the link
       current.next = prev
                     # Move prev to current
       prev = current
       current = next_node  # Move to next node
   return prev # New head of the reversed list
```

How to reverse a linked list - recursion

```
# Define the Node class
class Node:
    def __init__(self, data):
        self.data = data # Store data
        self.next = None # Pointer to the next node
# Function to reverse the linked list recursively
def reverseLinkedListRec(head):
    # Base case: if head is empty or only one node, it's already reversed
    if head is None or head.next is None:
        return head
    # Recursively reverse the rest of the list
    rest = reverseLinkedListRec(head.next)
    # Adjust the pointers
    head.next.next = head
    head.next = None
    # Return the new head of the reversed list
    return rest
```

Practice: check if there exists a cycle in a linked list (Optional)





Practice: check if there exists a cycle in a linked list (Optional)

```
# Define the Node class
class Node:
   def __init__(self, data):
        self.data = data # Store data
        self.next = None # Pointer to the next node
# Function to check if the linked list has a cycle
def hasCycle(head):
    slow = head # Tortoise moves one step at a time
   fast = head # Hare moves two steps at a time
   while fast and fast next:
       slow = slow.next # Move slow pointer by one
       fast = fast.next.next  # Move fast pointer by two
       if slow == fast:
           print("Cycle detected at node with data:", slow.data)
           return True
    print("No cycle detected.")
    return False
```

Circularly Linked List

- The tail of a linked list can use its next reference to point back to the head of the list
- Such a structure is usually called a circularly linked list



Example: Round-robin scheduler

- A round-robin scheduler iterates through a collection of elements in a circular fashion and "serves" each element by performing a given action on it
- Such a scheduler is used, for example, to fairly allocate a resource that must be shared by a collection of clients
- For instance, round-robin scheduling is often used to allocate slices of CPU time to various applications running concurrently on a computer

Implementing round-robin scheduler using standard queue

- A round-robin scheduler could be implemented with the standard queue, by repeatedly performing the following steps on queue Q:
 - 1) e = Q.dequeue()
 - 2) Service element e
 - 3) Q.enqueue(e)



Implement a Queue with a Circularly Linked List Class Node:

```
def __init__(self, element, pointer):
    self.element = element
    self.pointer = pointer
```

class CQueue:

```
def init (self):
    self. __tail = None
    self. size = 0
def __len__(self):
   return self. _____size
def is_empty(self):
   return self. __size == 0
def first(self):
    if self.is_empty():
        print('Queue is empty.')
    else:
        head = self. tail.pointer
        return head element
```

```
def dequeue(self):
    if self.is_empty():
        print('Queue is empty.')
    else:
        oldhead = self. __tail.pointer
        if self. __size == 1:
             self. tail = None
        else:
             self. __tail. pointer = oldhead. pointer
                                                       Skip the old head
        self. size -= 1
        return oldhead. element
def enqueue(self, e):
    newest = Node(e, None)
    if self.is empty():
                                                  A single self-pointed node
        newest.pointer = newest
    else:
        newest.pointer = self.__tail.pointer
        self. __tail. pointer = newest
                                                      Insert after the tail!
    self.__tail = newest
    self.___size += 1
```

Exercise: The Josephus Problem

• The Josephus Problem

- There are n people standing in a circle waiting to be executed. After the first man is executed, k - 1 people are skipped and the k-th man is executed. Then again, k-1 people are skipped and the k-th man is executed. The elimination proceeds around the circle (which is becoming smaller and smaller as the executed people are removed), until only the last man remains, who is given freedom.
- The task is to choose the place in the initial circle so that you survive, given n and k.



Doubly linked list

- For a singly linked list, we can efficiently insert a node at either end of a singly linked list, and can delete a node at the head of a list
- But we cannot efficiently delete a node at the tail of the list
- We can define a linked list in which each node keeps an explicit reference to the node before it and a reference to the node after it
- This kind of data structure is called doubly linked list

Head and tail sentinels

- In order to avoid some special cases when operating near the boundaries of a doubly linked list, it helps to add special nodes at both ends of the list: a header node at the beginning of the list, and a trailer node at the end of the list
- These "dummy" nodes are known as sentinels (or guards), and they do not store elements of the primary sequence



Inserting in the middle of a doubly linked list



Inserting at the head of the doubly linked list







Deleting from the doubly linked list



Code for the doubly linked list

```
class Node:
    def __init__(self, element, prev, nxt):
        self.element = element
        self.prev = prev
        self.nxt = nxt
```

class DLList:

```
def __init__(self):
    self.header = Node(None, None, None)
    self.trailer = Node(None, None, None)
    self.header.nxt = self.trailer
    self.trailer.prev = self.header
    self.size = 0
def __len__(self):
    return self.size
def is_empty(self):
    return self.size == 0
```

```
def insert_between(self, e, predecessor, successor): def main():
    newest = Node(e, predecessor, successor)
    predecessor.nxt = newest
    successor.prev = newest
    self.size+=1
    return newest

def delete_node(self, node):
    predecessor = node.prev

def delete_node(self, node):
    predecessor = node.prev

def delete_node(d.header.nxt.nxt)
```

```
d.iterate()
```

```
def iterate(self):
    pointer = self.header.nxt
    print('The elements in the list:')
    while pointer != self.trailer:
        print(pointer.element)
        pointer = pointer.nxt
```

node.prev = node.nxt = node.element = None

successor = node.nxt

element = node. element

self.size -=1

return element

predecessor.nxt = successor successor.prev = predecessor

Bubble sort

- Bubble sort is a simple sorting algorithm
- Its general procedure is:
- 1) Iterate over a list of numbers, compare every element i with the following element i+1, and swap them if i is larger
- 2) Iterate over the list again and repeat the procedure in step 1, but ignore the last element in the list
- 3) Continuously iterate over the list, but each time ignore one more element at the tail of the list, until there is only one element left



A longer example



Practice: Bubble sort over a standard list

```
def bubble(bubbleList):
    listLength = len(bubbleList)
    while listLength > 0:
        for i in range(listLength - 1):
            if bubbleList[i] > bubbleList[i+1]:
                buf = bubbleList[i]
                bubbleList[i]
                bubbleList[i] = bubbleList[i+1]
                bubbleList[i] = bubbleList[i+1]
                bubbleList[i+1] = buf
                listLength -= 1
                return bubbleList
```

```
def main():
    bubbleList = [3, 4, 1, 2, 5, 8, 0, 100, 17]
    print(bubble(bubbleList))
```

Practice: Bubble sort over a singly linked list





Solution:

```
from LinkedQueue import LinkedQueue
def LinkedBubble(q):
    listLength = q. size
    while listLength > 0:
        index = 0
        pointer = q. head
        while index < listLength-1:
            if pointer. element > pointer. pointer. element:
                buf = pointer.element
                pointer. element = pointer. pointer. element
  Swap values
                pointer.pointer.element = buf
            index += 1
            pointer = pointer.pointer
        listLength -= 1
    return q
```

```
def outputQ(q):
    pointer = q. head
    while pointer:
        print(pointer.element)
        pointer = pointer.pointer
def main():
    oldList = [9, 8, 6, 10, 45, 67, 21, 1]
    q = LinkedQueue()
    for i in oldList:
        q. enqueue(i)
    print ('Before the sorting...')
    outputQ(q)
    q = LinkedBubble(q)
    print()
    print('After the sorting...')
    outputQ(q)
```

Quick sort

- Quick sort is a widely used algorithm, which is more efficient than bubble sort
- The main procedure of quick sort algorithm is:
- 1) Pick an element, called a pivot, from the array
- 2) Partitioning: reorder the array so that all elements with values less than the pivot come before the pivot, while all elements with values greater than the pivot come after it (equal values can go either way). After this partitioning, the pivot is in its final position. This is called the partition operation
- 3) Recursively apply the above steps to the sub-array of elements with smaller values and separately to the sub-array of elements with greater values

Unsorted Array



A longer example



Practice: Quick sort over a standard list

```
def quickSort(L, low, high):
    i = low
    j = high
    if i \geq j:
        return L
    key = L[i]
    while i < j:
        while i \langle j  and L[j] \rangle= key:
             j = j-1
        L[i] = L[j]
        while i < j and L[i] \leq key:
            i = i+1
        L[j] = L[i]
    L[i] = key
    quickSort(L, low, i-1)
    quickSort(L, j+1, high)
    return L
```

Practice: Quick sort over a singly linked list

